



Success rate exponents for side-channel attacks

Sylvain Guilley, Annelie Heuser, Martial Ren, Olivier Rioul, Simon Sellem

► **To cite this version:**

Sylvain Guilley, Annelie Heuser, Martial Ren, Olivier Rioul, Simon Sellem. Success rate exponents for side-channel attacks. 16th Workshop on Cryptographic Hardware and Embedded Systems (CHES 2014), Sep 2014, Busan, South Korea. hal-02300002

HAL Id: hal-02300002

<https://hal.telecom-paris.fr/hal-02300002>

Submitted on 20 Aug 2021

HAL is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers.

L'archive ouverte pluridisciplinaire **HAL**, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d'enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.

Success Rate Exponents for Side-Channel Attacks

Sylvain GUILLEY
Annelie HEUSER
Martial REN
Olivier RIOUL
Simon SELLEM



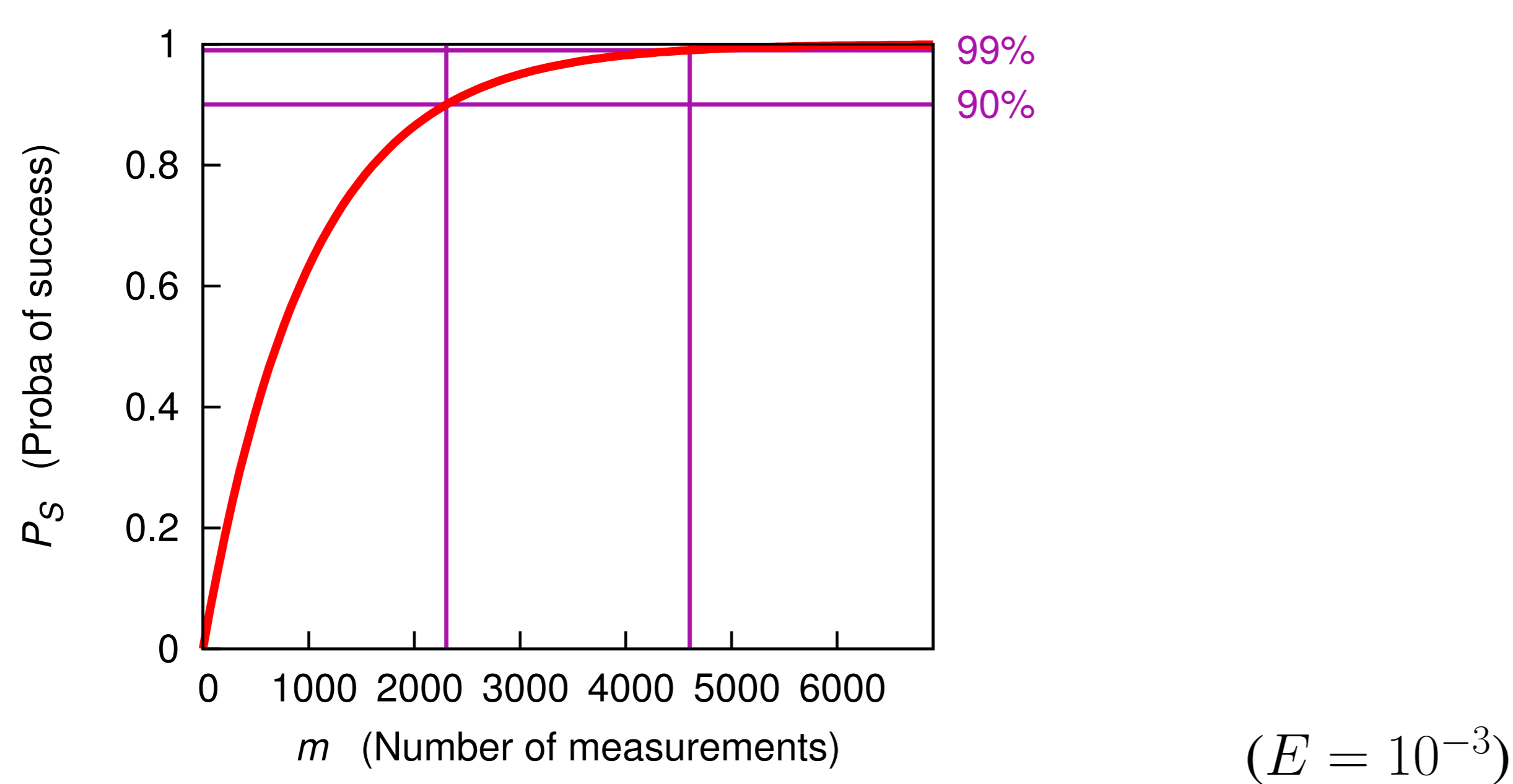
Goal

- Compute the exact probability of success $\mathbb{P}_S = \mathbb{P}(\hat{K} = K^*)$
- Rigorous mathematical computation of its first order exponent of success rate:

$$\mathbb{P}_S \approx 1 - e^{-mE} \quad \text{for some } E, \text{ where } m \text{ is the number of measurements.} \quad (1)$$

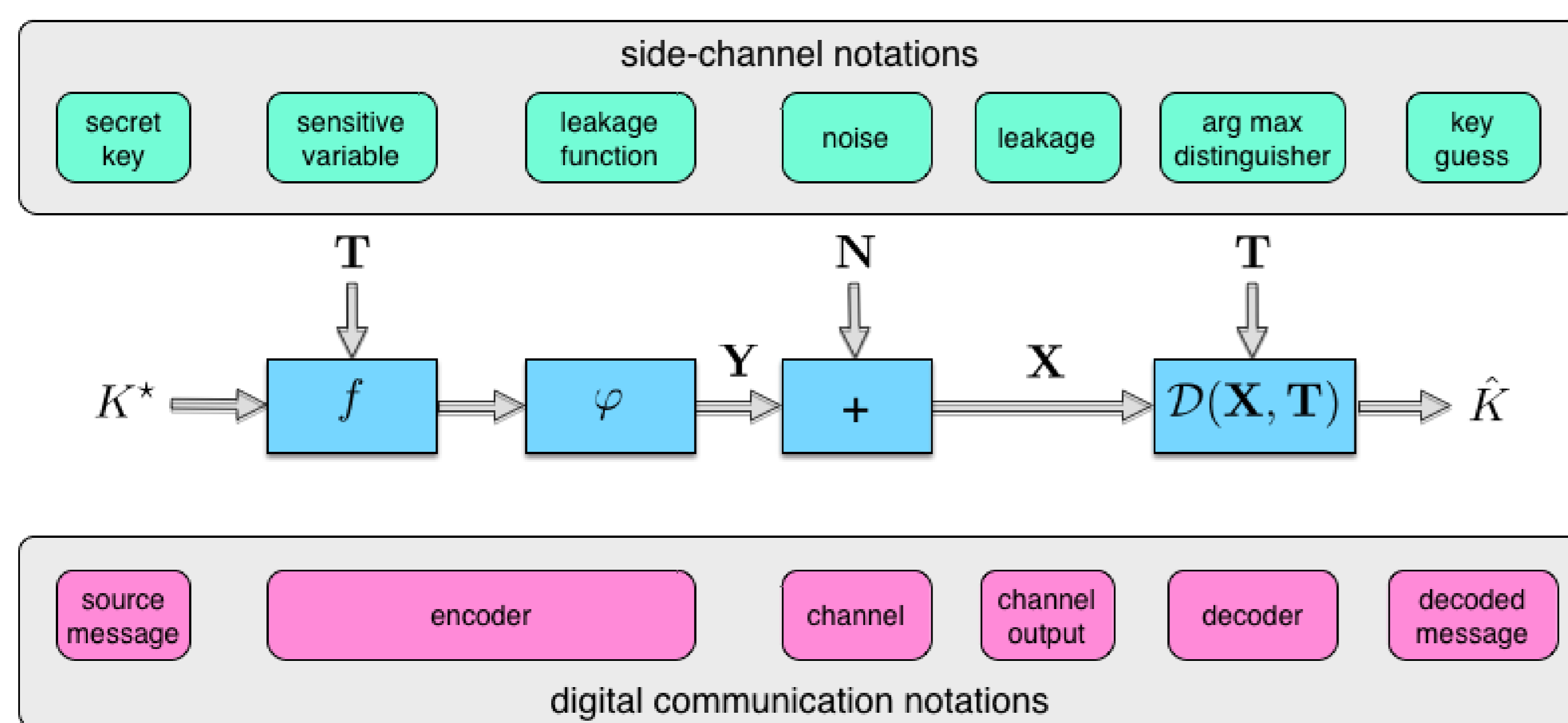
Example (with E independent of m)

- By Eq. (1), if $\mathbb{P}_S = 90\%$, then $m = \frac{\ln(10)}{E}$;
- Doubling the number of measurements $m \rightarrow 2m \Rightarrow \mathbb{P}_S = 99\%$.



Side-channel analysis

as a communication channel [HRG14]



Relationship with the state-of-the-art

At CRYPTO '99 [CJRR99], Chari, Jutla, Rao and Rohatgi gave a lower bound on the probability of success of mono-bit side-channel attacks. The bound is exponential in the number of queries.

Bibliographical references

- [CJRR99] Suresh Chari, Charanjit S. Jutla, Josyula R. Rao, and Pankaj Rohatgi. Towards Sound Approaches to Counteract Power-Analysis Attacks. In *CRYPTO*, volume 1666 of *LNCS*. Springer, August 15-19 1999. Santa Barbara, CA, USA. ISBN: 3-540-66347-9.
- [DZFL14] A. Adam Ding, Liwei Zhang, Yunsi Fei, and Pei Luo. A Statistical Model for Higher Order DPA on Masked Devices. Cryptology ePrint Archive, Report 2014/433, June 2014. <http://eprint.iacr.org/2014/433/> (to appear at CHES 2014).
- [FLD12] Yunsi Fei, Qiasi Luo, and A. Adam Ding. A Statistical Model for DPA with Novel Algorithmic Confusion Analysis. In Emmanuel Prouff and Patrick Schaumont, editors, *CHES*, volume 7428 of *LNCS*, pages 233–250. Springer, 2012.
- [HRG14] Annelie Heuser, Olivier Rioul, and Sylvain Guilley. Good is Not Good Enough — Deriving Optimal Distinguishers from Communication Theory. In Lejla Batina and Matt Robshaw, editors, *CHES*, Lecture Notes in Computer Science. Springer, 2014.
- [TPR13] Adrian Thillard, Emmanuel Prouff, and Thomas Roche. Success through Confidence: Evaluating the Effectiveness of a Side-Channel Attack. In Guido Bertoni and Jean-Sébastien Coron, editors, *CHES*, volume 8086 of *Lecture Notes in Computer Science*, pages 21–36. Springer, 2013.

Useful concepts

Definition 1 (First-Order Exponent Equivalence).

A sequence p_m of positive numbers admits a first-order exponent E_m if $\epsilon_m = E_m + \frac{1}{m} \ln p_m$ tends to zero as $m \rightarrow +\infty$. In this case we write:

$$p_m \approx e^{-mE_m}.$$

In practice, “ $m = 20$ ” is already “ $m \rightarrow +\infty$ ”.

Result for Gaussian noise

When $Y = \alpha X(k) + N$, with $N \sim \mathcal{N}(0, \sigma^2)$ is the noise:

$$E = \frac{1}{8\sigma^2} \min_{k \neq k'} \mathbb{E}(Y(k) - Y(k'))^2 \quad (2)$$

$$= \frac{1}{2} \times \text{SNR} \times \min_{k \neq k'} \kappa_{k,k'}, \quad (3)$$

- where $\text{SNR} = \frac{\alpha^2}{\sigma^2}$, and
- where $\kappa_{k,k'} = \frac{1 - \rho(Y(k), Y(k'))}{2}$ is the confusion coefficient [FLD12].

Proof (for the optimal distinguisher when the noise is Gaussian, i.e., the opposite norm-2)

$$D_k^{\text{opt}}(\mathbf{x}, \mathbf{t}) = -\|\mathbf{x} - \mathbf{y}(k)\|_2. \quad (4)$$

Nota bene: bold letters \mathbf{x} represent series: $\mathbf{x} = (x_i)_{1 \leq i \leq m}$.

We define the *marginal pairwise probability of failure*:

$$\mathbb{P}_{k \rightarrow k' | \mathbf{t}} = \mathbb{P}(\|\mathbf{X} - \mathbf{Y}(k)\| \leq \|\mathbf{X} - \mathbf{Y}(k')\| | \mathbf{T} = \mathbf{t}). \quad (5)$$

Lemma 1 (Q -expression of the Pairwise Failure Probability).

Let $Q(x) = \frac{1}{2} \text{erf}\left(\frac{x}{\sqrt{2}}\right) = \frac{1}{\sqrt{2\pi}} \int_x^{+\infty} e^{-t^2/2} dt$ be the complementary cdf of a standard normal $\mathcal{N}(0, 1)$. We have:

$$\mathbb{P}_{k \rightarrow k' | \mathbf{t}} = Q\left(\frac{\|\mathbf{y}(k) - \mathbf{y}(k')\|}{2\sigma}\right).$$

Lemma 2 (first-order Gaussian Q exponent).

As $x \rightarrow +\infty$, $Q(x) \lesssim e^{-x^2/2}$, in the sense that $Q(x) \leq e^{-x^2/2}$ for all $x \geq 0$ and $\frac{\ln Q(x)}{x^2} \rightarrow -\frac{1}{2}$.

Dominated convergence theorem: $\mathbb{E}[\ln \mathbb{P}_{k \rightarrow k' | \mathbf{T}} / m] \rightarrow -\frac{\mathbb{E}(Y(k) - Y(k'))^2}{8\sigma^2}$, when $m \rightarrow +\infty$. Combining yields the result. \square

Our results generalize the closed-form expressions of the success probability recently obtained in the case of normal noise when correlation is used as a distinguisher [FLD12, TPR13, DZFL14]. First order exponents are an alternative to the *relative distinguishing margin* (RDM [WO11]) because:

- they are completely related to the ideal criterion (namely, the success rate) for large (even moderately large) number of measurements m ,
- they can be computed explicitly for many distinguishers to indicate how performance relates to noise power, model, etc. without resorting to experimentation.